

Question 1: What is the output of the following program?

```
#include <iostream>
using namespace std;

int s(int n)
{
    if (n == 0) return 0;
    return n + s(n-1);
}

int main()
{
    cout << s(10) << endl;
}
```

Question 2: What is the running time of $s(n)$?

(In these questions, use the carat symbol (\wedge) for exponentiation). **Question 3:** What is the output of the following program?

```
#include <iostream>
using namespace std;

string t(string v)
{
    string tmp;
    if (v.size() == 0) return "";
    tmp.push_back(v[0]);
    return t(v.substr(1)) + tmp;
}

int main()
{
    cout << t("Fred") << endl;
}
```

Question 4: What is the running time of $t(v)$, where $n = v.size()$?

Question 5: What would be the running time of $t(v)$ if v were a reference parameter?

Answers to the Clicker Questions

Question 1

- $s(10)$ returns $10 + s(9)$.
- $s(9)$ returns $9 + s(8)$.
- $s(8)$ returns $8 + s(7)$.
- And so on
- $s(0)$ returns 0

So, $s(n)$ returns $n + (n-1) + \dots + 2 + 1$, which is $n(n+1)/2$. The answer is 55.

Question 2

Answer: It makes n recursive calls, and each call does $O(1)$ work (besides the recursion). So the answer is $O(n)$.

Question 3

- $s(\text{"Fred"})$ sets **tmp** to "F" and calls $s(\text{"red"})$. It returns $s(\text{"red"}) + \text{"F"}$.
- $s(\text{"red"})$ sets **tmp** to "r" and calls $s(\text{"ed"})$. It returns $s(\text{"ed"}) + \text{"r"}$.
- $s(\text{"ed"})$ sets **tmp** to "e" and calls $s(\text{"d"})$. It returns $s(\text{"d"}) + \text{"e"}$.
- $s(\text{"d"})$ sets **tmp** to "d" and calls $s(\text{""})$. It returns $s(\text{""}) + \text{"d"}$.
- $s(\text{""})$ returns "".
- So, $s(\text{"d"})$ returns "d".
- So, $s(\text{"ed"})$ returns "de".
- So, $s(\text{"red"})$ returns "der".
- So, $s(\text{"Fred"})$ returns "derF".

In other words it reverses the string: "derF".

Question 4

Answer: It makes n recursive calls; However each recursive call creates a substring of size $(n-1)$ and then copies it when calling $s()$ on the substring. So the running time of this is $n + (n-1) + (n-2) + \dots + 2 + 1$, which is $O(n^2)$.

Question 5

With the reference parameter, a copy of the substring is no longer made. However, making the substring still takes $O(n)$ work, so the running time is still $O(n + (n-1) + (n-2) + \dots + 2 + 1)$, which is still $O(n^2)$.